### High-Performance Computational Electromagnetic Modeling Using Low-Cost Parallel Computers

July 14, 1997

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### Beowulf System (Hyglac)

◆ 16 PentiumPro PCs, each with 2.5 Gbyte disk, 128 Mbyte memory, Fast Ethernet card.

◆ Connected using 100Base-T network, through a 16-way crossbar switch.

◆ Theoretical peak performance: 3.2 GFlop/s.

 ◆ Achieved sustained performance: 1.26 GFlop/s.



### Hyglac Cost

◆ Hardware cost: \$54, 200 (9/18/96)

- 16 (CPU, disk, memory)
- − 1 (16-way crossbar, monitor, keyboard, mouse)
- ◆ Software cost: \$0
  - Public Domain OS, compilers, tools, libraries.
- ◆ 256 PE T3D: \$8 million (1/94)
  - including all software and maintenance for 3 years.
- ◆ 16 PE T3E: \$1 million
  - 3 to 4 times more powerful then T3D.



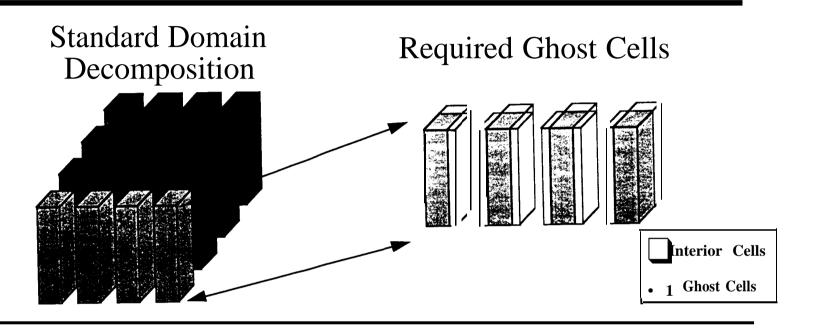
### Hyglac Performance (vs. T3D)

Γ	Hyglac	T3D	T3D
	(MPI)	, (MPI)	(shmem)
CPU Speed (MHz)	200	150	150
Peak Rate (MFlop/s)	200	150	150
L1, L2 Cache Size	8i+8d, 256	8i+8d, 0	8i+8d, 0
(Kbytes)		:	à
Memory Bandwidth	0.36?	1.4	1.4
(Gbit/s)			.3
Communic ation	150	35	1.8
Latency (µs)			
Communication	66	225	280-970
Bandwidth (Mbit/s)			

◆ Next, examine performance of EMCC FDTD and PHOEBUS (FE) codes...



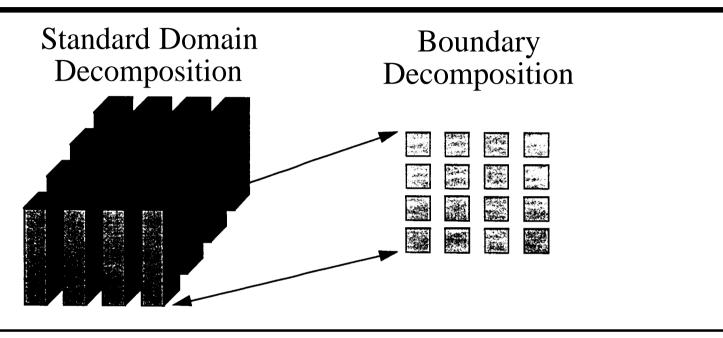
### FDTD Interior Communication



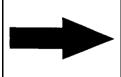
One plane of ghost cells must be communicated to each neighboring processor each time step.



### FDTD Boundary Communication



Extra work required along faces (boundary conditions, wave source, far-field data locus).



Different decomposition required for good load balance.

Data at 4 faces must be redistributed twice each time step!!



### FDTD Timing

<u> </u>	T3D (shmem)	T3 (MP		
Interior	1.8	1.8	1.1	1.1
Computation	(1.3*)		(1.3*)	2 0
Interior	0.00	· • • • • • • • • • • • • • • • • • • •		J0
Communication.		·		
Boundary	0.19	0.19	0.14	0.42
Computation				
Boundary	0.04	1.5	50.1	0.0
Communication				-
Total	2.0	3.5	55.1	5.5
	$(1.5^*)$	$(3.0^*)$	*	

("using assembler kernel)

All timing data is in CPU seconds/simulated time step, for a global grid size of 282 x 362 x 102, distributed on 16 processors.



### FDTD Timing, cont.

- **◆** Computation:
  - Hyglac is 35-65% faster than T3D.
- **♦** Communication:
  - T3D: MPI is 4 to 9 times slower than shmem.
  - Hyglac MPI is 30-50 times slower than T3D MPI.
- ◆ Good (or even acceptable) performance requires rewriting/modifying code.



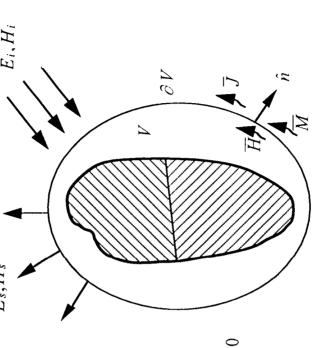
# PHOEBUS Coupled Formulation

 $\overline{E} = \overline{E}_s + \overline{E}_i$   $\overline{H} = \overline{H}_s + \overline{H}_i$ 

For three unknowns,

$$\overline{H}$$
,  $\overline{J}$ ,  $\overline{M}$ 

The following three equations must be solved:



$$\frac{1}{-j\omega\varepsilon_0} \int_{\nu} \left[ \left( \nabla \times \overline{T} \right) \frac{1}{\varepsilon_r} \left( \nabla \times \overline{H} \right) - k_0^2 \overline{T} \ \mu_r \overline{H} \right] dV + \left[ \overline{T} \cdot \overline{M} \ dS = 0 \right]$$
(finite element equation)

 $\int_{\partial V} \hat{n} \times \overline{U} \cdot \left[ \hat{n} \times \overline{H} - \overline{J} \right] dS = 0$  (essential boundary condition)

$$Z_{e}[\overline{M}] + Z_{h}[\overline{J}] = V_{i}$$
 (combined field integra equation)



### PHOEBUS Coupled Equations

$$\begin{bmatrix} K & C & 0 \\ C^{\dagger} & 0 & Z_0 \\ 0 & Z_M & Z_J \end{bmatrix} \begin{bmatrix} H \\ M \\ J \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ V_{inc} \end{bmatrix}$$

- ◆ This matrix problem is filled and solved by PHOEBUS.
  - The K submatrix is a sparse finite element matrix.
  - The Z submatrices are integral equation matrices.
  - The C submatrices are coupling matrices between the FE and IE matrices.



### PHOEBUS Two Step Method

$$\begin{bmatrix} K & C & 0 \\ C^{\dagger} & 0 & Z_0 \\ 0 & Z_M & Z_J \end{bmatrix} \begin{bmatrix} H \\ M \\ J \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ V \end{bmatrix}$$

$$H = -K^{-1}CM$$

$$\begin{bmatrix} -C^{\dagger}K^{-1}C & Z_{0} \\ Z_{M} & Z \end{bmatrix} \begin{bmatrix} M \\ J \end{bmatrix} \begin{bmatrix} 0 \\ \Psi \end{bmatrix}$$

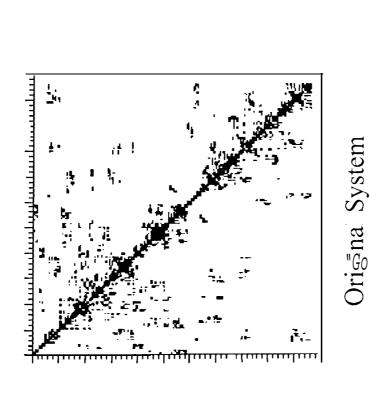
- ♦ Find  $-C^{\dagger}K^{-1}C$  using QMR on each row of C, building x rows of  $K^{-1}C$ , and multiplying with  $C^{\dagger}$ .
- ◆ Solve reduced system as a dense matrix.
- If required, save  $K^{-1}C$  to solve for H.

### PHOEBUS Decomposition

- ◆ Matrix Decomposition
  - Assemble complete matrix.
  - Reorder to equalize row bandwidth.
    - » Gibbs-Poole-Stockmeyer
    - )) SPARSPAK' s GENRCM
  - partition matrix in slabs or blocks.
  - Each processor receives slab of matrix elements.
  - Solve matrix equation.



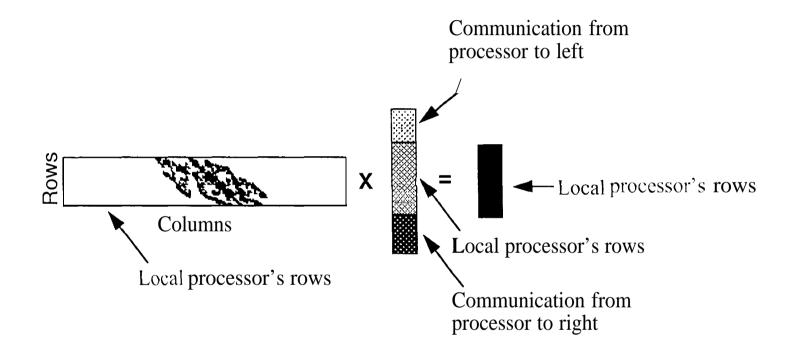
### PHOEBUS Matrix Reordering



for Minimum Bandwidth System after Reordering

Using SPARSPAK's GENRCM Reordering Routine

### PHOEBUS Matrix-Vector Multiply





### PHOEBUS Solver Timing

Model: dielectric cylinder with 43,791 edges, radius = 1 cm, height = 10 cm, permittivity = 4.0, at 5.0 GHz

	T3D T3D H (shmem) (MPI)	lyglac MPI)
Matrix-vector Multiply Computation	1290   1290	590
Matrix-Vector Multiply Communication		3260
Other Work Total	407 415 1800 1980	1360 5220

Time of Convergence (CPU seconds), solving using pseudo-block QMR algorithm for 116 right hand sides.



### • Computation:

- Hyglac CPU works 55% faster than T3D CPU.
- For sparse arithmetic, large secondary cache provides substantial performance gain.



### **♦** Communication:

- T3D MPI is 2.5 times slower than T3D shmem.
- Hyglac MPI is 10 times slower than T3D MPI.
- The code was rewritten to use manual packing and unpacking, to combine 10 small messages into 1 medium-large message.
- On another "supercomputer", this led to a factor of 10 speed-up in communications.



### ♦ Other work:

- Mostly BLAS 1 -type operations (dot, norm, scale, vector sum, copy)
  - » Heavily dependent on memory bandwidth.
- Also, some global sums
  - )) Over all PEs.
  - )) Length: number of RHS/block complex words.
- 3.5 times slower than T3D.



### Conclusions

- ◆ If message sizes are not too small, adequate communication performance is possible.
- + Computation rates are good, provided there is adequate data reuse.
- ◆ Low-cost parallel computers can provide better performance/\$ than traditional parallel supercomputers, for limited problem sizes, and with some code rewriting.



### Conclusions, cont.

- ◆ Both EMCC FDTD code and PHOEBUS QMR solver should scale to larger numbers of processors easily and well, since most communication is to neighbors.
- ◆ User experience has been mostly good:
  - Lack of both support and documentation can be frustrating.
  - Easy to use, once you know how use it.



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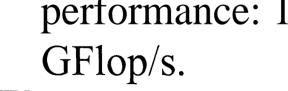
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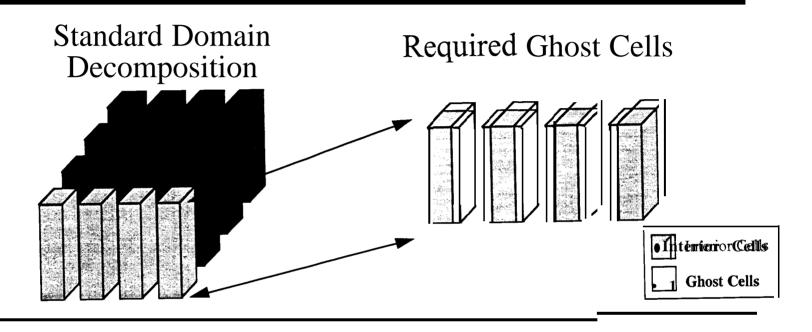
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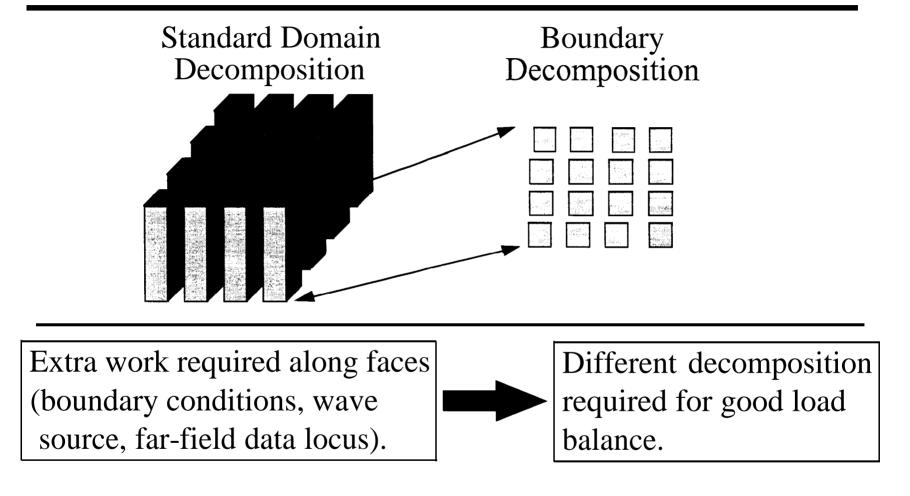
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### FDTD Timing

	T3D T3D	Hyglac	Hyglac
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	· · · · · · ·	Good Load	Bad Load
	:	Balance)	Balance)
Interior	1.8 1.8	1.1	1.1
Computation	(1.3") $(1.3")$		
Interior	$0.007 \mid 0.08$	3.8	3.8
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Boundary	0.19 0.19	0.14	0.42
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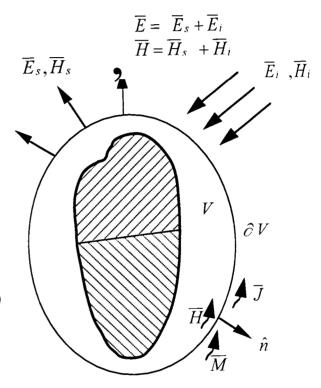
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$$\begin{vmatrix} K & C & 0 & | H \\ C^{\dagger} & 0 & Z_0 & | M \\ O & Z_M & Z_J & | J \end{vmatrix} = \begin{vmatrix} 0 \\ 0 \\ V_{inc} \end{vmatrix}$$

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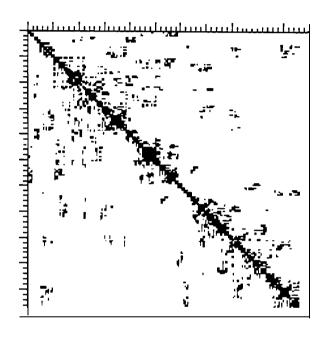


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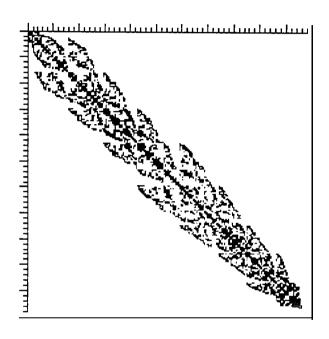
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### PHOEBUS Matrix Reordering



Original System

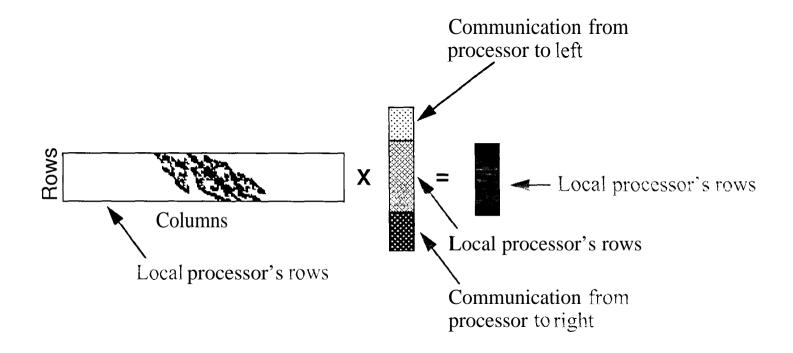


System after Reordering for Minimum Bandwidth

Using SPARSPAK's GENRCM Reordering Routine



### PHOEBUS Matrix-Vector Multiply





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Other Work	407 415 1360
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